

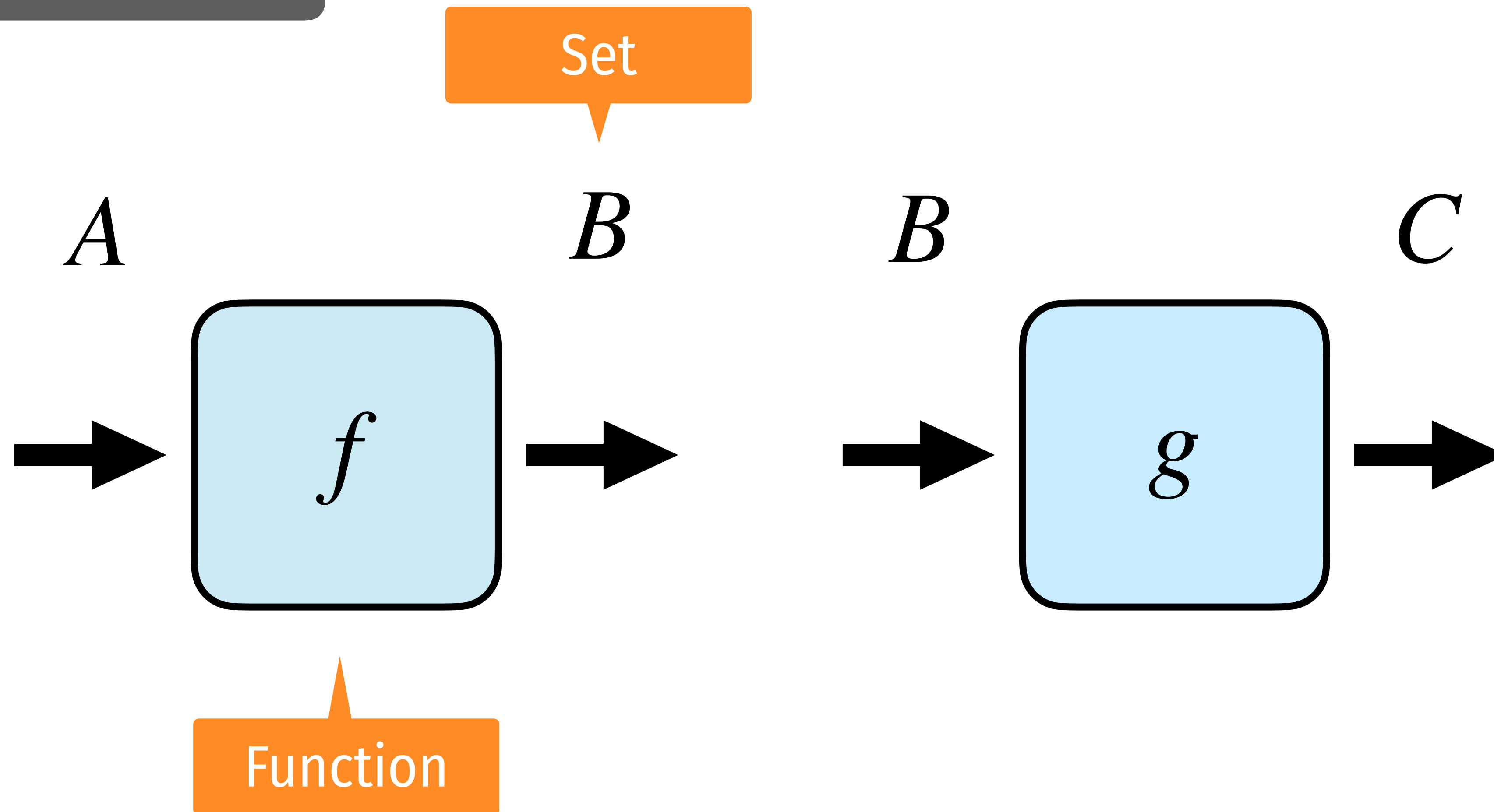
The Selection Monad Transformer and Higher-Order Games

**Winter Meeting in Lisbon - Mathematics & Applications
18/12/2023**

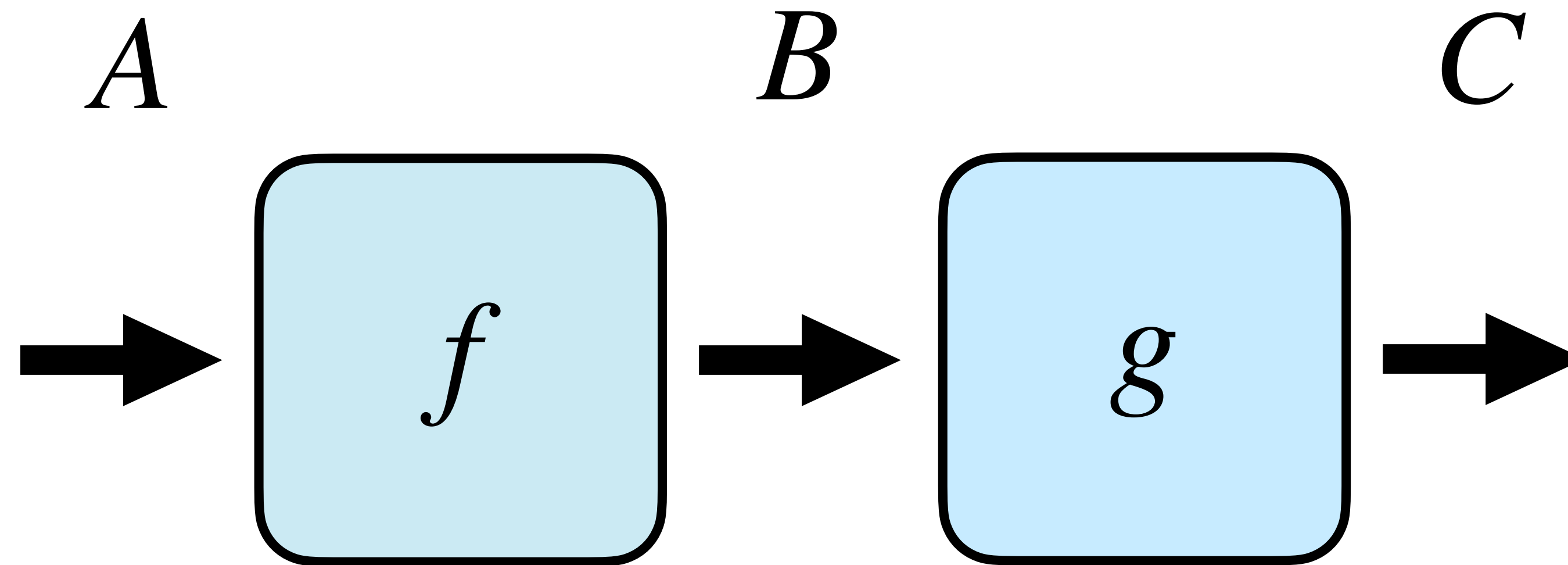
**Paulo Oliva
(jww Martin Escardo)
Queen Mary University of London**

Modularity and Composition

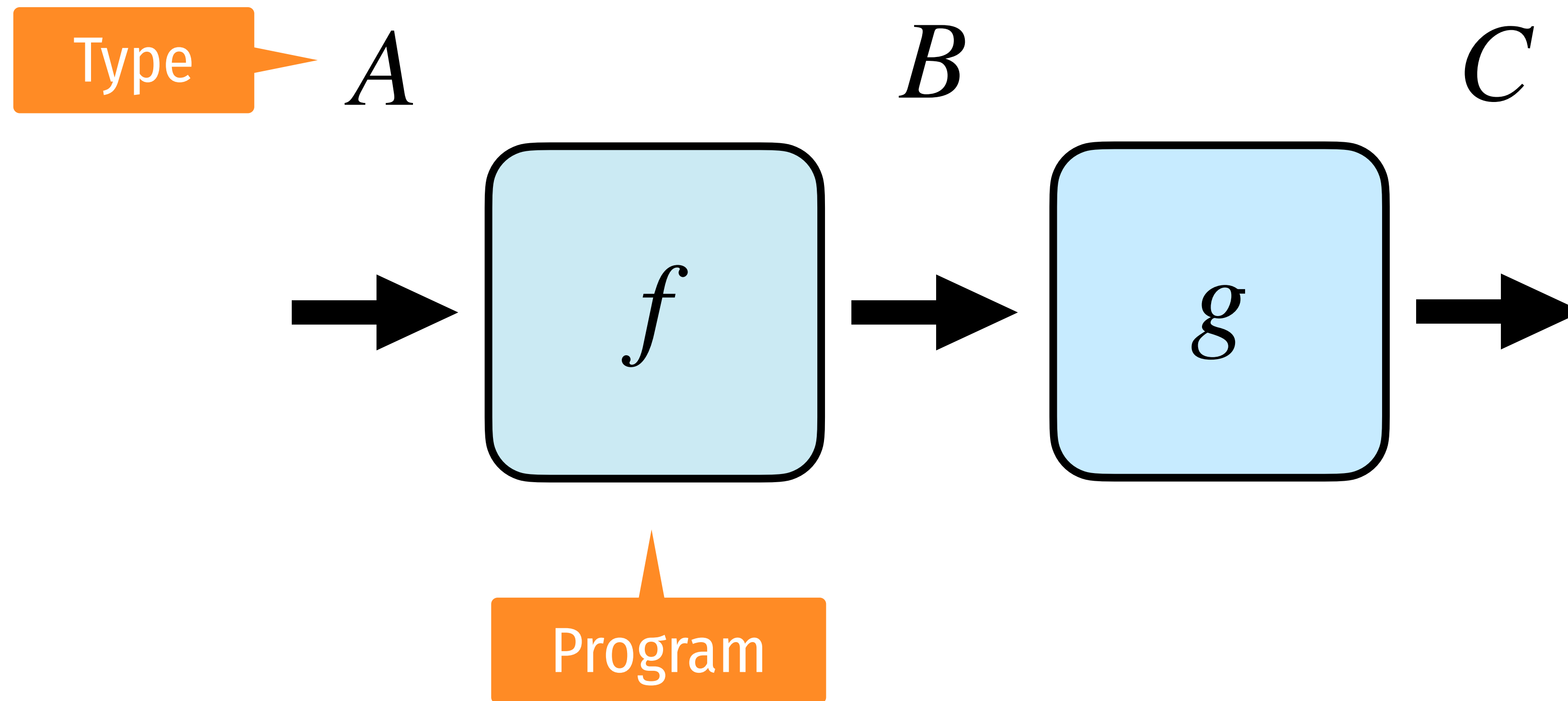
Composition



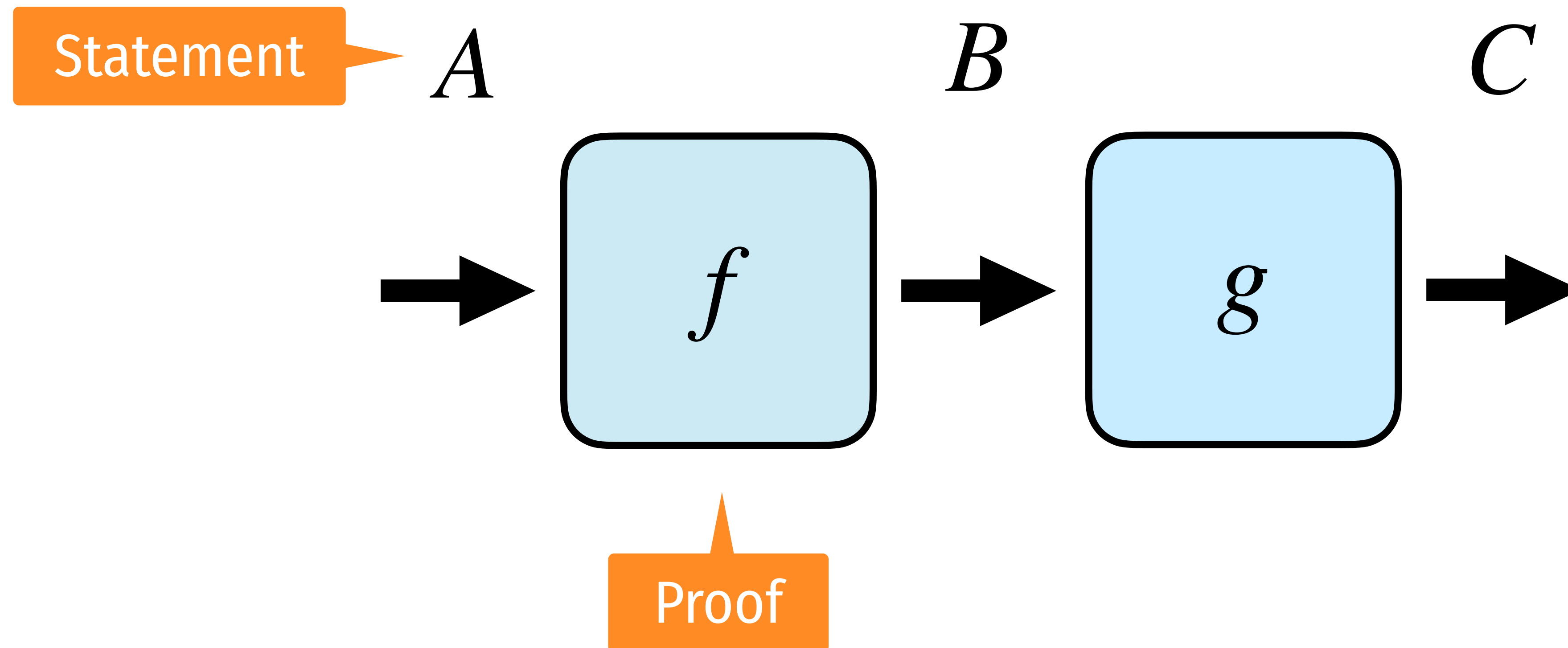
Composition



Composition of Programs

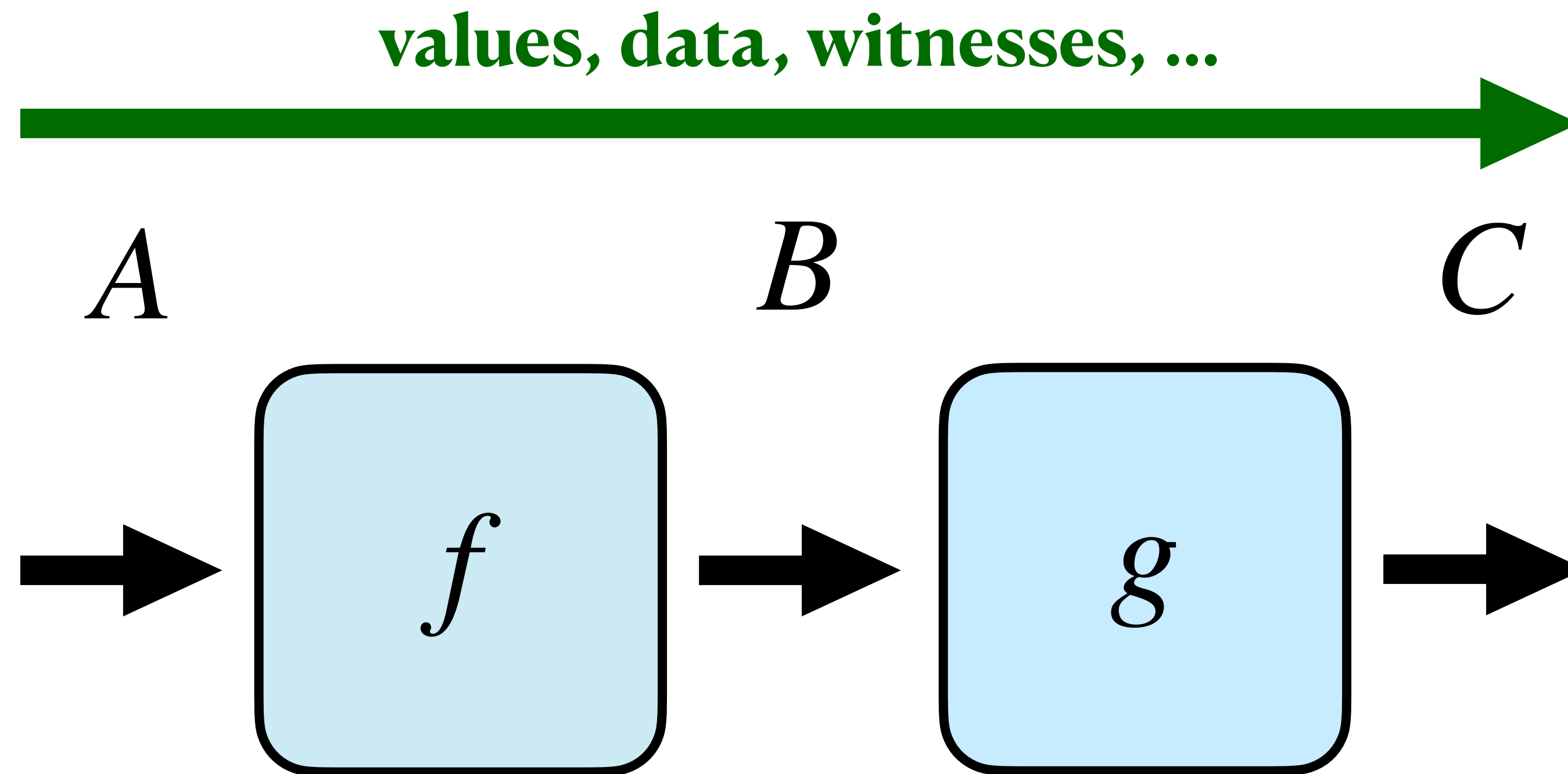


Composition of Proofs

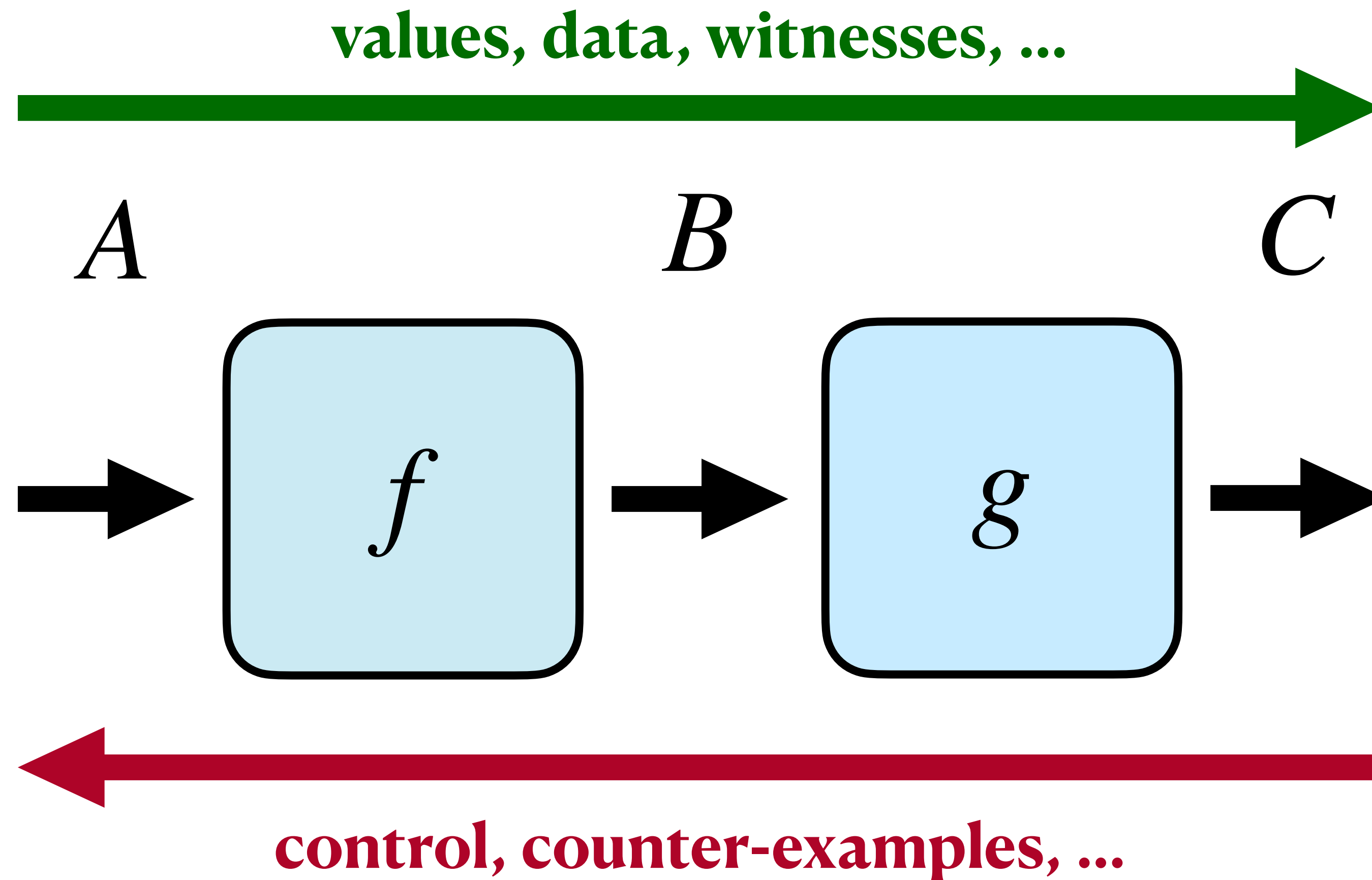


Flow and Counter-Flow

Flow



Counter-Flow



ÜBER EINE BISHER NOCH NICHT BENÜTZTE ERWEITERUNG DES FINITEN STANDPUNKTES

von Kurt GÖDEL, Princeton

P. Bernays hat wiederholt darauf hingewiesen ¹, dass angesichts der Tatsache der Unbeweisbarkeit der Widerspruchsfreiheit eines Systems mit geringeren Beweismitteln als denen des Systems selbst eine Überschreitung des Rahmens der im Hilbertschen Formalismus der Mathematik nötig ist, um die Widerspruchsfreiheit der Mathematik, ja sogar um die der klassischen Zahlentheorie zu beweisen. Da die finite Mathematik als die der *analysis* definiert ist ², so bedeutet das (wie auch von Bernays in *Revue internationale de philosophie*, 34 (1935), p. 62 und 69, exp. wurde), dass man für den Widerspruchsfreiheitsbeweis gewisse *abstrakte* Begriffe braucht. Dabei



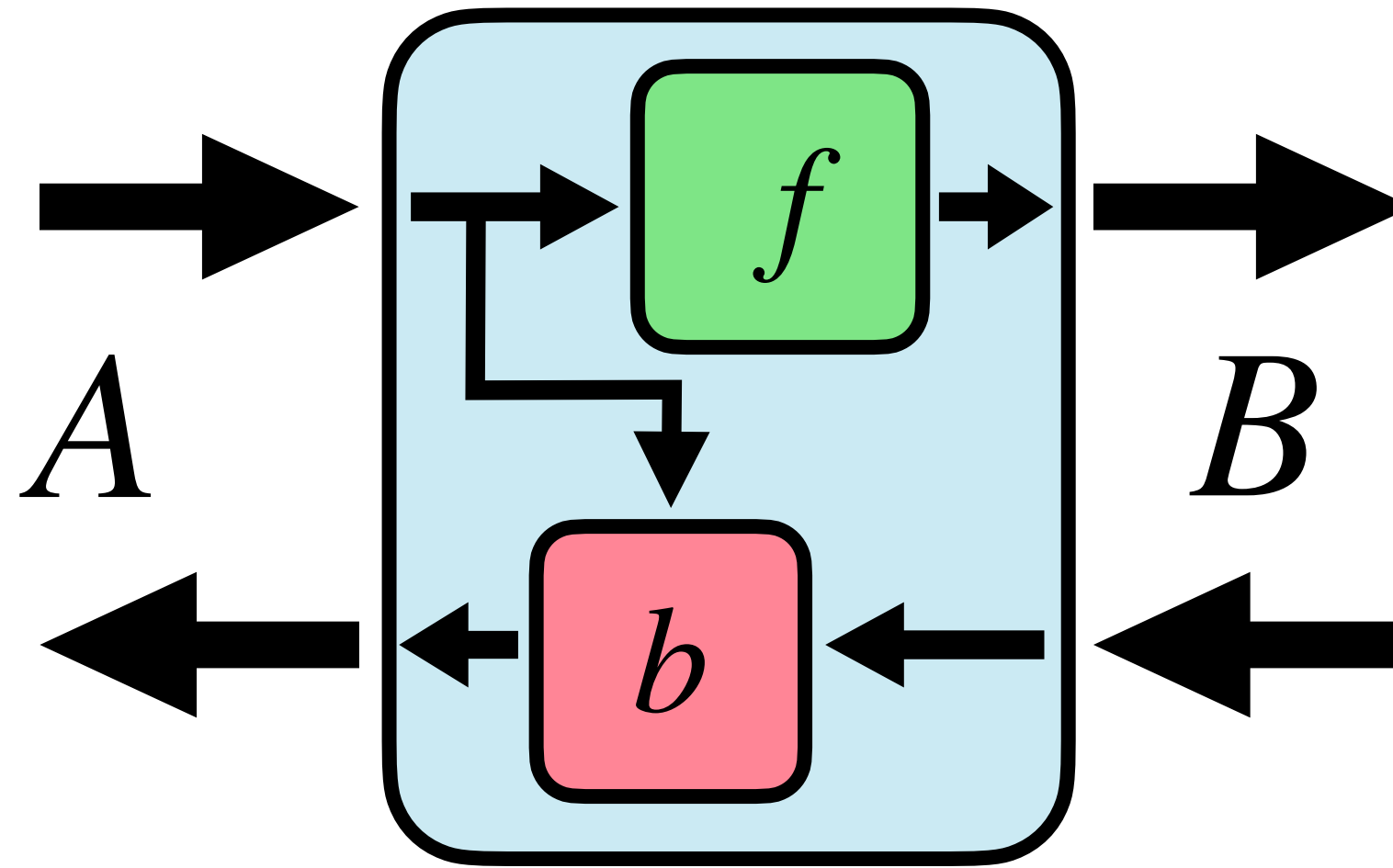
Volume 12, Issue 3-4

December 1958

Pages 280-287

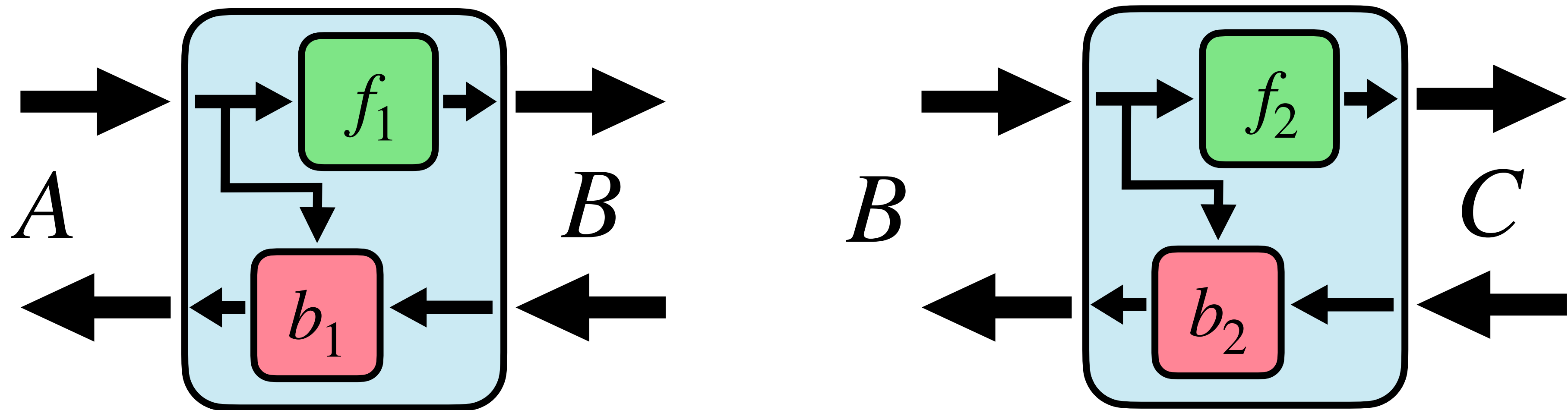
Flow and Counter-Flow

Relation $A \subseteq A^+ \times A^-$

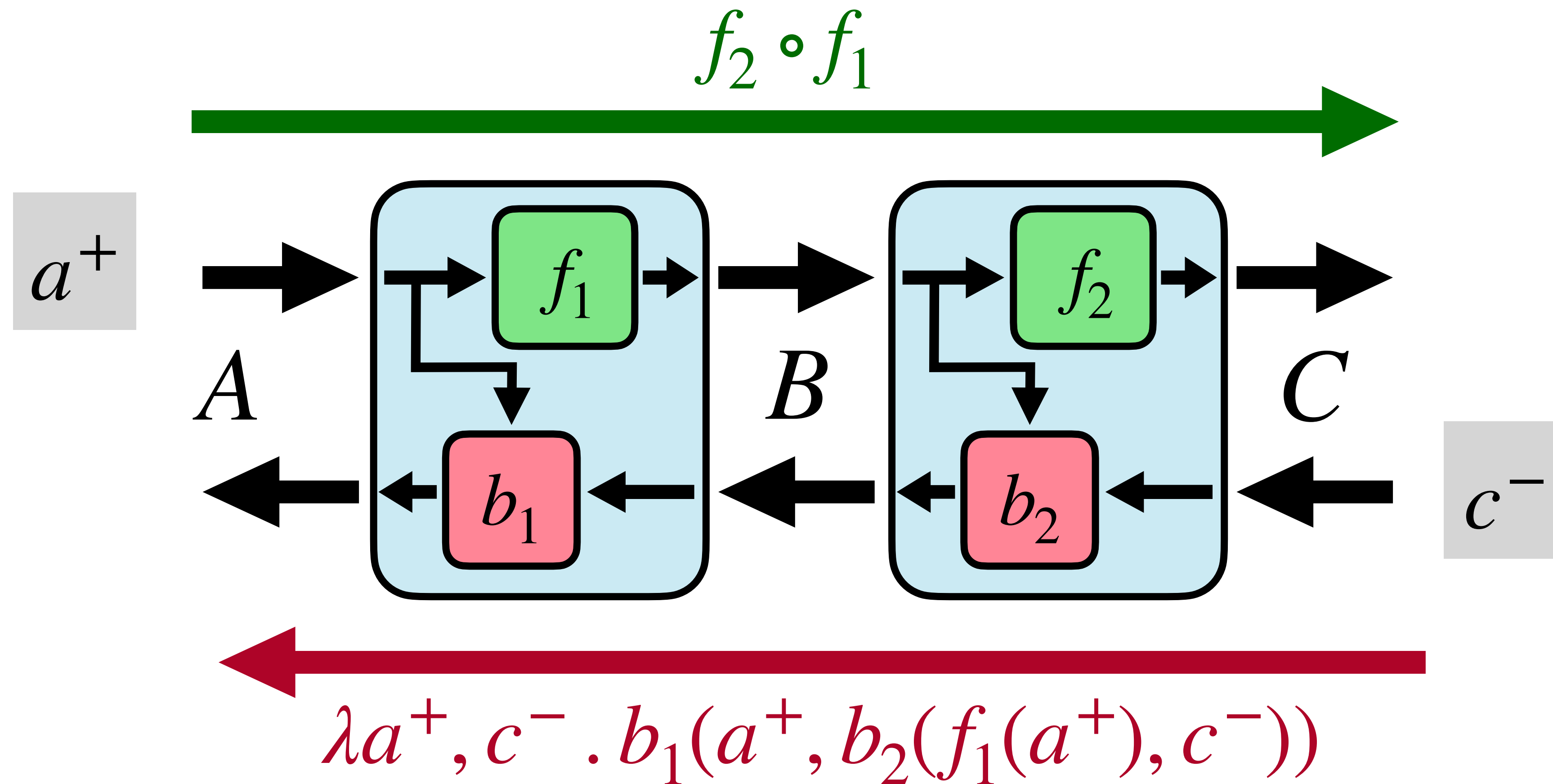


Pair of functions $(f, b): (A^+ \rightarrow B^+) \times (A^+ \times B^- \rightarrow A^-)$

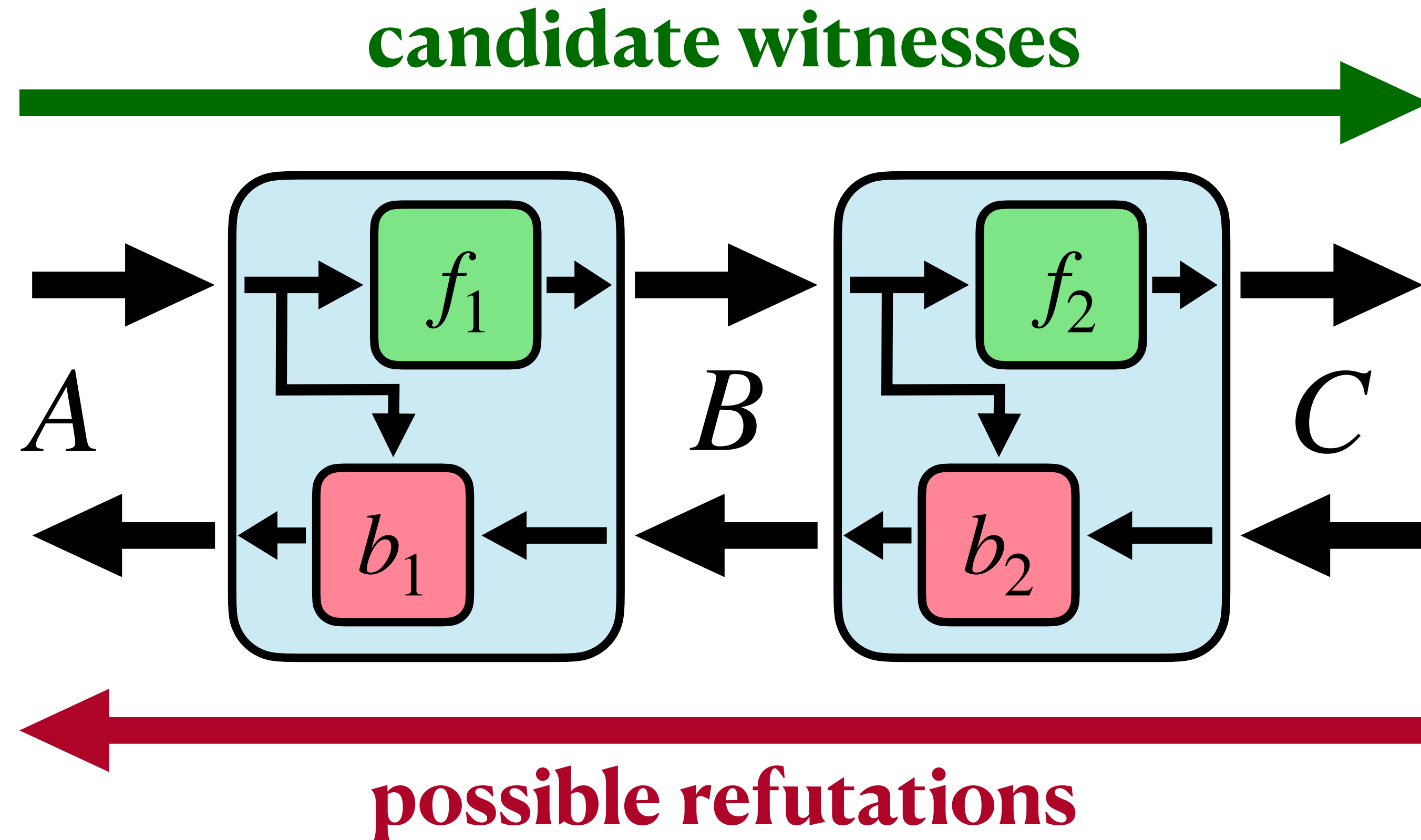
Composition



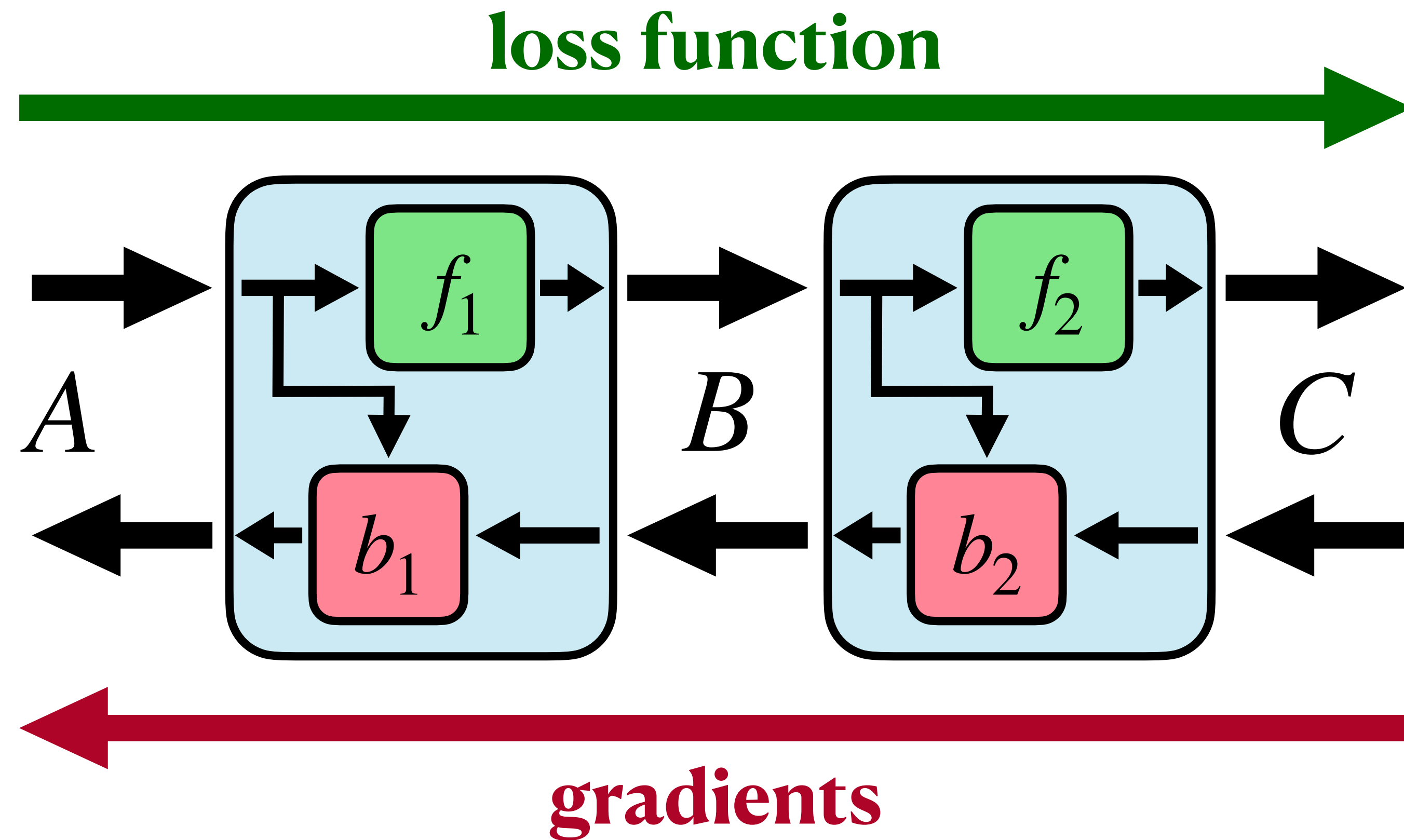
Composition



Gödel's Dialectica Interpretation



Neural Networks and Back Propagation



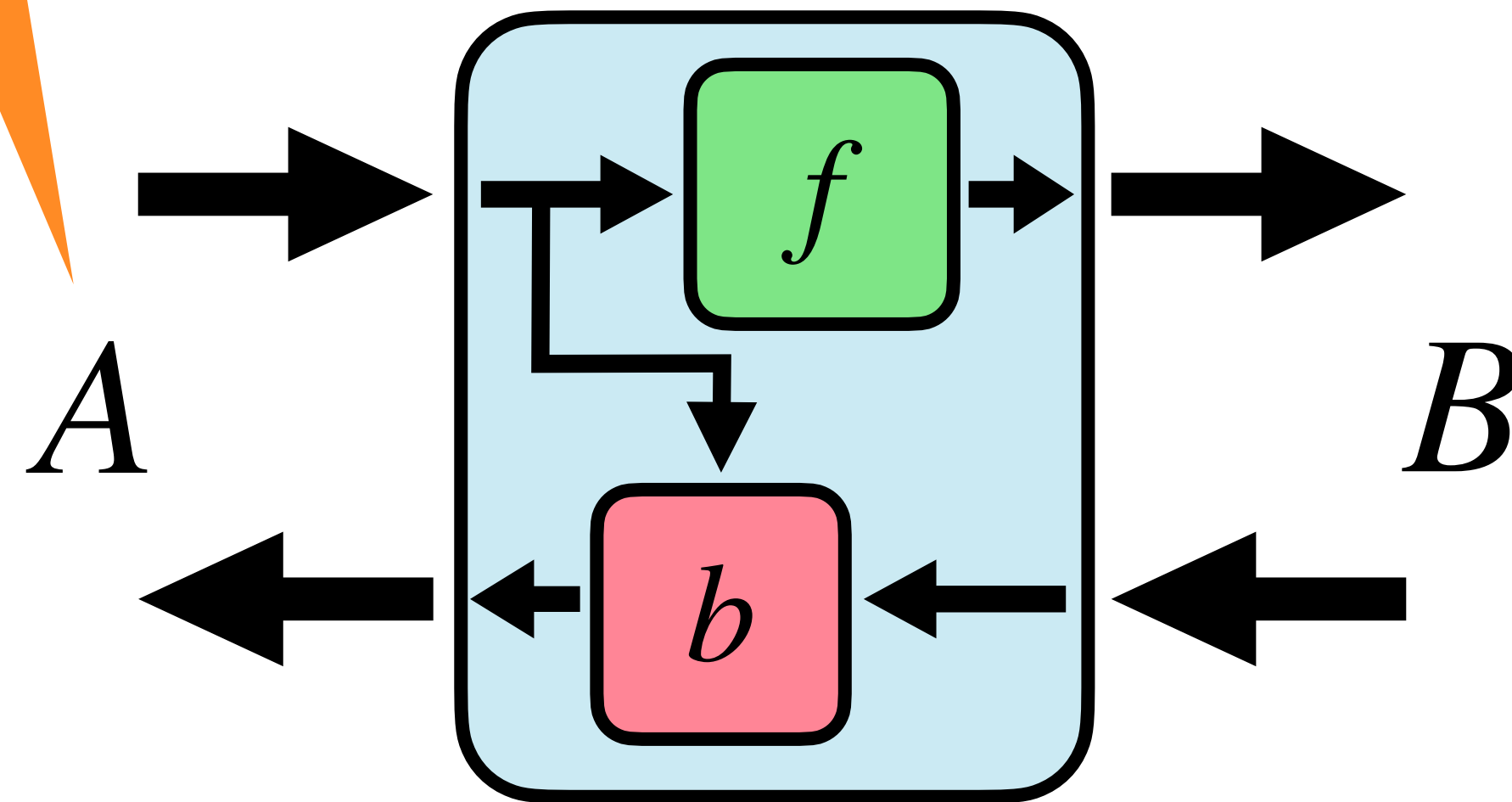
Games

Games

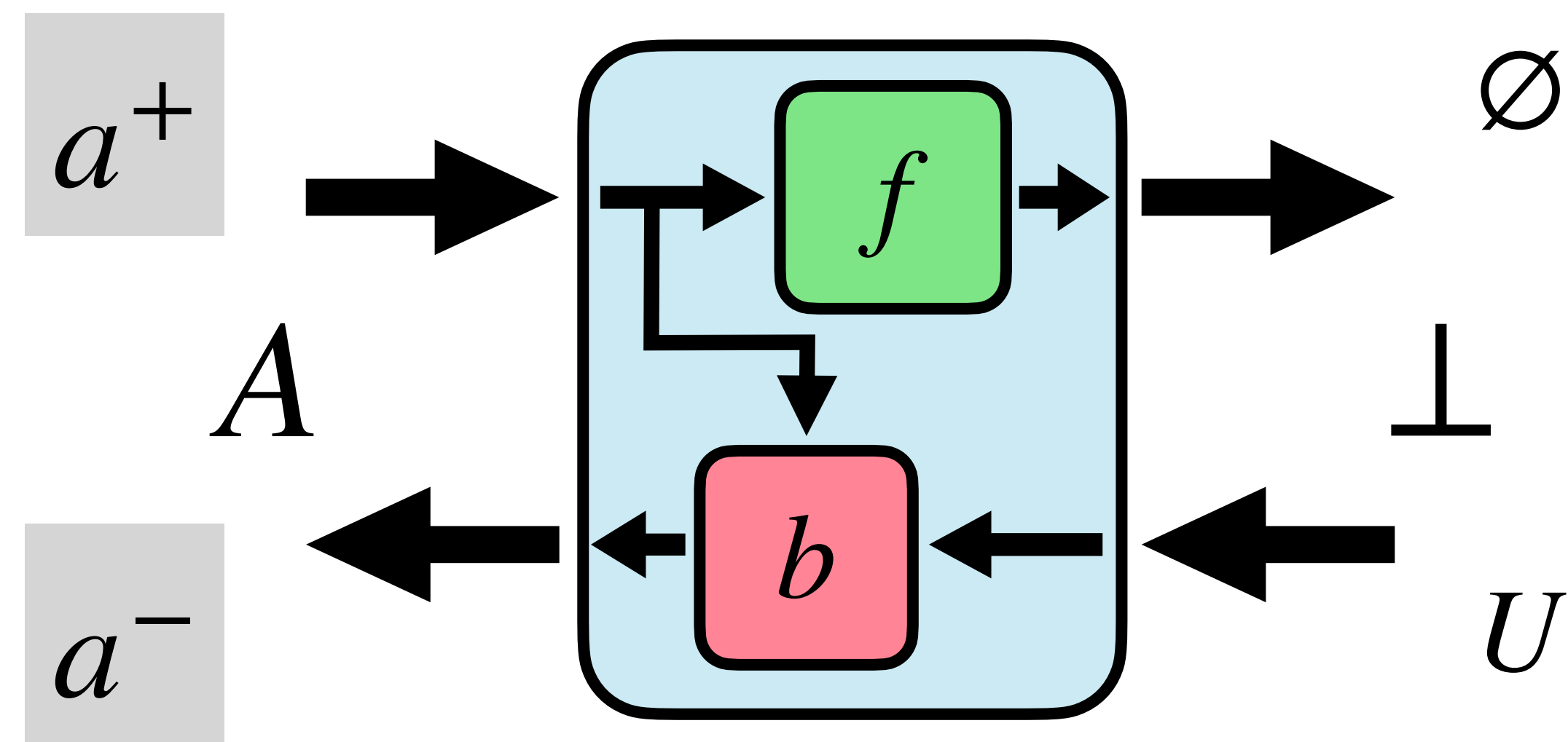
Play game B relative to game A

Relation $A \subseteq A^+ \times A^-$

Game
 $A^+ = \text{Eloise}$
 $A^- = \text{Abelard}$

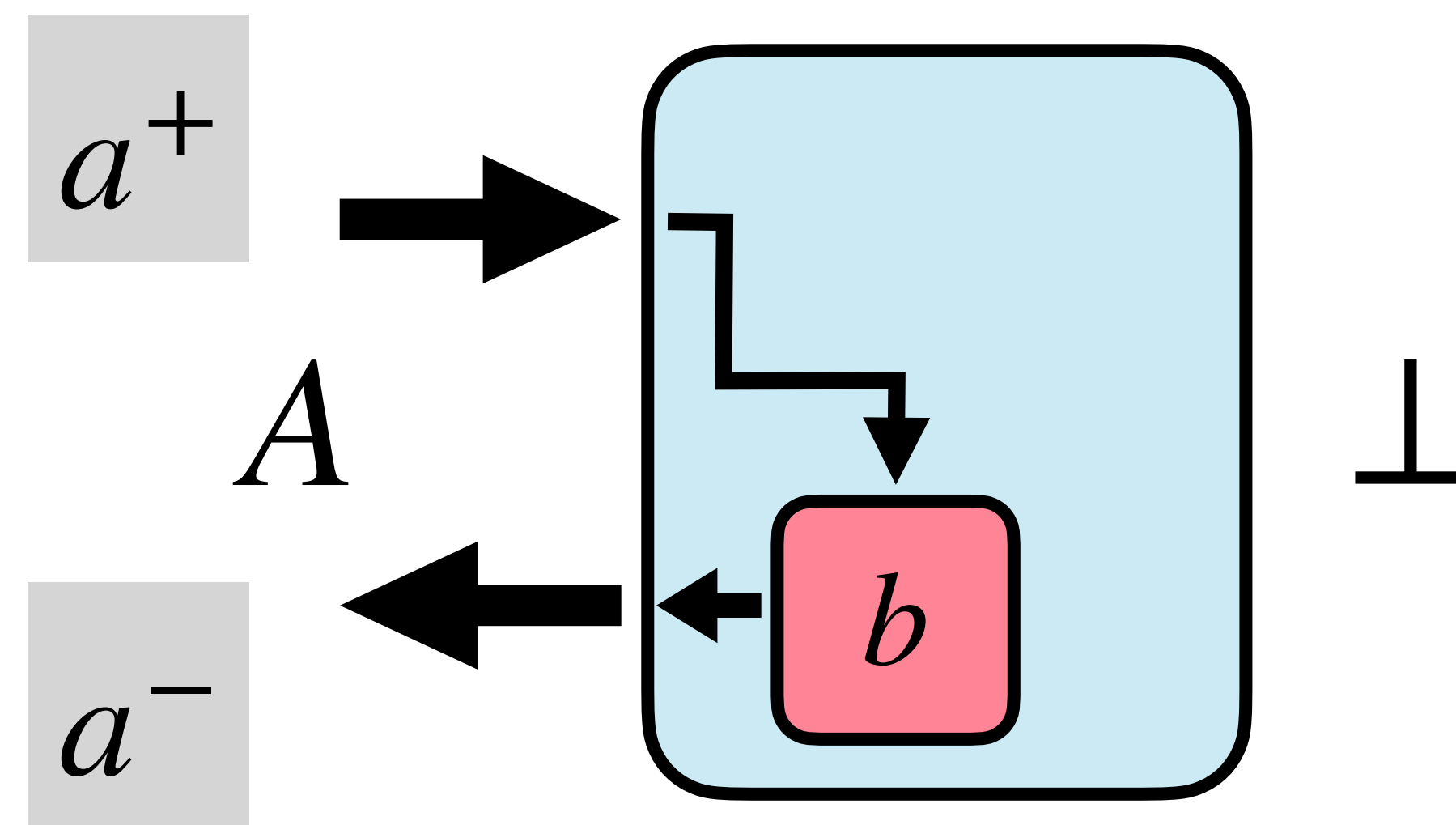


The $\neg A$ Game



The $\neg A$ Game

$$\text{Game } \neg A \subseteq (A^+ \rightarrow A^-) \times A^+$$



The $\neg\neg A$ Game

$$\text{Game } \neg\neg A \subseteq ((A^+ \rightarrow A^-) \rightarrow A^+) \times (A^+ \rightarrow A^-)$$

$$(X \rightarrow R) \rightarrow X$$

selection functions

Selection Functions and Quantifiers

$$\varepsilon: (X \rightarrow R) \rightarrow X$$

selection function

$$\varphi: (X \rightarrow R) \rightarrow R$$

quantifier

$$\varphi(p) = p(\varepsilon(p))$$

attainability

Selection Functions and Quantifiers

$$\varepsilon: (X \rightarrow \mathbb{R}) \rightarrow X$$

selection function

$$\exists: (X \rightarrow \mathbb{B}) \rightarrow \mathbb{B}$$

quantifier

$$\exists(p) = p(\varepsilon(p))$$

attainability

Selection Functions and Quantifiers

$$\operatorname{argmax} : (X \rightarrow \mathbb{R}) \rightarrow X$$

selection function

$$\max : (X \rightarrow \mathbb{R}) \rightarrow \mathbb{R}$$

quantifier

$$\max(p) = p(\operatorname{argmax}(p))$$

attainability

Selection and Continuation Monad

$$J_R X = (X \rightarrow R) \rightarrow X$$

selection monad

$$K_R X = (X \rightarrow R) \rightarrow R$$

continuation monad

$$\eta_X^K: X \rightarrow K_R X$$

$$\eta_X^J: X \rightarrow J_R X$$

$$\beta_X^K: K_R X \rightarrow (X \rightarrow K_R Y) \rightarrow K_R Y$$

$$\beta_X^J: J_R X \rightarrow (X \rightarrow J_R Y) \rightarrow J_R Y$$

Tensor of Quantifiers and Selection Functions

We can compose
selection functions
and quantifiers!

$$\otimes : K_R X \times K_R Y \rightarrow K_R(X \times Y)$$

$$\otimes : J_R X \times J_R Y \rightarrow J_R(X \times Y)$$

Higher-Order Games

Higher-Order Game

 R

outcomes

 X_i

possible moves
at round i

 $q: \prod_i X_i \rightarrow R$

outcome function
(outcome given a play)

 $\varphi_i: (X_i \rightarrow R) \rightarrow R$

quantifiers
(preferred outcome given context)

REVIEW

Sequential games and optimal strategies

BY MARTÍN ESCARDÓ¹ AND PAULO OLIVA^{2,*}

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This article gives an overview of recent work on the theory of sequential games. It explains the intuition behind these higher type objects, sequential game whose optimal strategies can be computed by backward selection functions. Several instances of this game are considered, such as fixed point theory, topology, game theory, higher type games. These examples are intended to illustrate how the fundamental results of game theory can be used to solve problems in other fields. Strategies based on products of selection functions permit the construction of solutions to problems in game theory.

Keywords: selection function; Nash equilibrium; backward induction; bar recursion; Tychonoff theorem



julesh
@_julesh_

The paper Sequential Games and Optimal Strategies by Martín Escardó and Paulo Oliva, which single-handedly got me into game theory and convinced me to apply to do a PhD with Paulo
royalsocietypublishing.org/doi/full/10.1098/rspa.2010.0471



Paul Graham ✓ @paulg · Sep 14, 2019

If you had to pick one thing you've read that changed the course of your life, what would it be? (ht @rivatez for the question.)

9:04 PM · Sep 14, 2019

Higher-Order Game

 R

outcomes

 X_i

moves

 $q: \prod_i X_i \rightarrow R$

outcome function

 $\varphi_i: (X_i \rightarrow R) \rightarrow R$

goal functions

 $\left(\bigotimes \varphi_i \right) (q)$

optimal outcome

 $\left(\bigotimes \varepsilon_i \right) (q)$

optimal strategy

Monadic Strategies

Selection Monad Transformer

$$J_R X = (X \rightarrow R) \rightarrow X$$

selection monad

Given any monad TX such that R is a T -algebra ($\alpha: TR \rightarrow R$)...

$$J_R^T X = (X \rightarrow R) \rightarrow TX$$

... also a monad

Recent & Current Work

- Recent work:
 - Use of **powerset monad** for interpretation of countable choice via Herbrand interpretation
 - Extension of the framework to deal with **dependent types**
- Current work:
 - **Probability monad**: Sub-optimal or irrational players
 - **Reader monad**: Alpha-beta pruning

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- 📄 M. Escardó and P. Oliva, **Selection functions, bar recursion and backward induction**, *Mathematical Structures in Computer Science*, 20(2):127-168, 2010
- 📄 M. Escardó and P. Oliva, **Sequential games and optimal strategies**, *Proc. of the Royal Society A*, 467:1519-1545, 2011
- 📄 M. Escardó and P. Oliva, **The Herbrand functional interpretation of the double negation shift**, *The Journal of Symbolic Logic*, 82(2):590-607, 2017
- 📄 M. Escardó and P. Oliva, **Higher-order games with dependent types**, *Theoretical Computer Science*, 974, 2023